

DYCK PATHS OF KNIGHT MOVES*

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The NE (north-east) and SE (south-east) steps of Dyck paths are replaced by NNE, NEE, SSE and SEE steps (i.e. moves of a chess knight from left to right) to obtain "Dyck paths of knight moves". In order to study these paths (called A -paths), we need to introduce several closely related families of paths and to find a system of equations between the corresponding generating functions. The generating function $A = A(t)$ for A -paths satisfies a fourth-degree polynomial equation which can be solved. The same method also works for the general (s, r) -knight but the number of new families of paths to be introduced in the system grows rapidly ($4r - 2$ equations for an (s, r) -knight). The generating function for these paths still satisfies a polynomial equation (for example of degree 8 when $(s, r) = (1, 3)$ or $(2, 3)$). One may think of all these paths as words in the letters x, y, \bar{y} ; the corresponding languages are algebraic.

Introduction

Dyck paths are counted by Catalan numbers and are in bijection with Dyck words which form a nice algebraic language on two letters (see [3, 6, 10, 11]). In the present text, we replace NE (north-east) and SE (south-east) steps by NNE, NEE, SSE and SEE steps (i.e. moves of a chess knight from left to right). In order to enumerate these paths (called A -paths), we need to introduce two closely related families of paths, find a system of equations relating generating functions and solve it. The generating function $A = A(t)$ for A -paths satisfies a fourth-degree equation (Theorem 2.13) while for Dyck path it is a well-known second-degree equation (Theorem 1.9). The same method also works for the study of the case of a general (s, r) -knight. More new paths have to be introduced and the corresponding generating functions satisfy an intricate system of $4r - 2$ equations (Theorem 3.5). The generating function for A -paths still satisfies a polynomial equation (for example of degree 8 when $(s, r) = (1, 3)$ or $(2, 3)$). One may think of all these paths as words in the letters x, y, \bar{y} ; the corresponding languages are algebraic.

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1. Classical Dyck paths

Definition 1.1. A *Dyck path of length n* is a path in $\mathbb{Z} \times \mathbb{Z}$ which

- (i) is made only of steps NE (north-east) or SE (south-east),
- (ii) starts at $(0, 0)$ and ends at $(n, 0)$, and
- (iii) never goes strictly below the x -axis.

Remark 1.2. It is well known (see [2, 3, 6, 9–11]) that Dyck paths are in bijection with *Dyck words, ballot sequences, well-formed sequences of parentheses, 2-lines standard tableaux, binary trees, ordered trees*, and all these are counted by Catalan numbers.

Definition 1.3. Let d_n be the number of these paths and $d_n^{(k)}$ for $k \geq 0$ be the number of those which hit (not counting at $(0, 0)$) the x -axis exactly k times. Note that $d_0^{(0)} = 1$ and $d_n^{(0)} = 0$ for $n > 0$; moreover $d_n = 0$ when n is odd.

Definition 1.4. Let the corresponding generating functions be

$$D = \sum_{n \geq 0} d_n t^n \quad \text{and} \quad D^{(k)} = \sum_{n \geq 0} d_n^{(k)} t^n.$$

Lemma 1.5. We have $D = \sum_{k \geq 0} D^{(k)}$.

Lemma 1.6. We have $D^{(k)} = (D^{(1)})^k$ for $k \geq 1$.

Proof. Obviously, $d_n^{(k)} = \sum_{m=1}^{n-1} d_m^{(1)} d_{n-m}^{(k-1)}$, where in this sum m corresponds to the first place where the path hits, at $(m, 0)$, the x -axis. This implies $D^{(k)} = D^{(1)} \cdot D^{(k-1)}$ and the result. \square

Lemma 1.7. We have $D = (1 - D^{(1)})^{-1}$.

Lemma 1.8. We have $D^{(1)} = t^2 D$.

Proof. Let $n \geq 2$. Removing the first and last steps of a Dyck path which hits the x -axis only at $(0, 0)$ and $(n, 0)$, we get bijectively an arbitrary Dyck path of length $n - 2$. This proves $d_n^{(1)} = d_{n-2}$ which implies the result. \square

Theorem 1.9. The generating function D for Dyck paths satisfies the equation $t^2 D^2 - D + 1 = 0$.

Corollary 1.10. We have $D = \frac{1}{2} (1 - (1 - 4t^2)^{1/2}) / t^2$.

Corollary 1.11. We have $d_{2m} = \binom{2m}{m} / (m + 1)$ (the m th-Catalan number).

2. Dyck paths of knight moves

Definition 2.1. A Dyck path of knight moves of size n (which we call an *A-path*) is a path in $\mathbb{Z} \times \mathbb{Z}$ which

- (i) is made only of steps NNE, NEE, SSE and SEE,
- (ii) starts at $(0,0)$ and ends at $(n,0)$,
- (iii) never goes strictly below the x -axis (see Fig. 1).

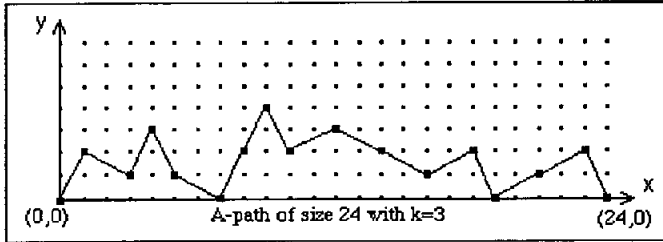


Fig. 1.

Definition 2.2. Let a_n be the number of these paths and $a_n^{(k)}$ the number of those which hit the x -axis (except at $(0,0)$) exactly k times. Note that $a_0^{(0)} = 1$ and $a_n^{(0)} = 0$ for $n > 0$.

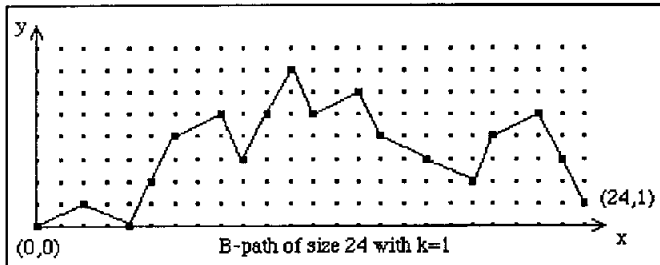


Fig. 2.

Definition 2.3. *B-paths* of size n satisfy the three conditions of Definition 2.1 except that they end at $(n,1)$ (see Fig. 2). Let b_n and $b_n^{(k)}$ be defined as before. Note that $b_0^{(0)} = b_1^{(0)} = 0$ but $b_n^{(0)} > 0$ for $n \geq 2$.

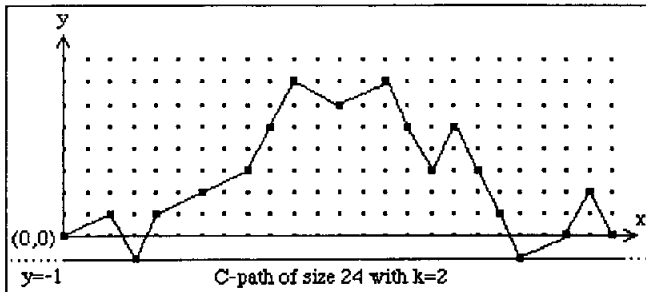


Fig. 3.

Definition 2.4. *C*-paths of size n satisfy the three conditions of Definition 2.1 except that they are allowed to go below the x -axis but always remain above the line $y = -1$ (see Fig.3). Let c_n and $c_n^{(k)}$ be as before except that $k, k \geq 0$, is now the number of times the *C*-path hits the line $y = -1$.

Let $A, A^{(k)}, B, B^{(k)}, C$ and $C^{(k)}$, be the corresponding generating functions.

Lemma 2.5. *We have*

- (a) $A = \sum_{k \geq 0} A^{(k)}$,
- (b) $B = \sum_{k \geq 0} B^{(k)}$,
- (c) $C = \sum_{k \geq 0} C^{(k)}$.

Lemma 2.6. *We have for $k \geq 1$*

- (a) $A^{(k)} = A^{(1)} \cdot A^{(k-1)} = (A^{(1)})^k$,
- (b) $B^{(k)} = A^{(1)} \cdot B^{(k-1)} = (A^{(1)})^k \cdot B^{(0)}$,
- (c) $C^{(k)} = B^{(0)} \cdot B^{(k-1)} = (B^{(0)})^2 \cdot (A^{(1)})^{k-1}$.

Proof. For example, a look at Fig.3 should convince the reader of the validity of the formula, $c_n^{(k)} = \sum_{m=1}^{n-1} b_m^{(0)} \cdot b_{n-m}^{(k-1)}$, which implies (c). The other identities are proved similarly. \square

Lemma 2.7. *We have $A^{(1)} = t^4 A + 2t^3 B + t^2 C$.*

Proof. In an arbitrary $A^{(1)}$ -path (i.e. which hits the x -axis only at $(0,0)$ and $(n,0)$) erase the first and last steps:

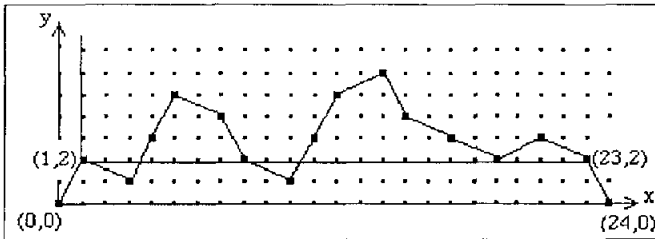


Fig. 4.

Case 1. First step is NNE and last step is SSE. Since one gets an arbitrary “*C*-path” from $(1,2)$ to $(n-1,2)$ there are c_{n-2} such paths (see Fig. 4).

Case 2. First step is NEE and last step is SEE. Since one gets an arbitrary “*A*-path” from $(2,1)$ to $(n-2,1)$ there are a_{n-4} such paths.

Case 3. First step is NEE and last step is SSE. Since one gets an arbitrary “*B*-path” from $(2,1)$ to $(n-1,2)$ there are b_{n-3} such paths.

Case 4. First step is NNE and last step is SEE; by symmetry there are b_{n-3} such paths. This is a bijective proof of $a_n^{(1)} = c_{n-2} + a_{n-4} + 2b_{n-3}$, which implies the result. \square

Lemma 2.8. *We have $B^{(0)} = t^2A + tB$.*

Proof. This time erasing the first step of a B -path which hits the x -axis only at $(0, 0)$ gives a bijective proof of the formula $b_n^{(0)} = a_{n-2} + b_{n-1}$, which implies the result. \square

Lemma 2.9. *We have $C^{(0)} = A$.*

Proof. The result follows by definition. \square

Lemma 2.10. *We have $A = (1 - A^{(1)})^{-1}$.*

Proof. The result follows from Lemmas 2.5(a) and 2.6(a). An A -path is a chain of $A^{(1)}$ -paths. \square

Lemma 2.11. *We have $B = t^2A^2(1 - tA)^{-1}$.*

Proof. $B = \sum_{k \geq 0} B^{(k)} = B^{(0)} \cdot \sum_{k \geq 0} A^{(k)} = B^{(0)} \cdot A = (t^2A + tB) \cdot A = t^2A^2 + tAB$. This implies the result. \square

Lemma 2.12. *We have $C = t^2AB + tB^2 + A$.*

Proof. $C = \sum_{k \geq 0} C^{(k)} = \sum_{k \geq 1} C^{(k)} + A = \sum_{k \geq 1} B^{(0)} \cdot B^{(k-1)} + A = B^{(0)} \cdot B + A = t^2AB + tB^2 + A$. \square

Theorem 2.13. *The generating function $A = A(t)$ for A -paths of size n satisfies the equation*

$$t^4A^4 - (2t^3 + t^2)A^3 + (t^4 + 2t^2 + 2t)A^2 - (2t + 1)A + 1 = 0.$$

Proof. By Lemma 2.7, $A^{(1)} = t^4A + 2t^3B + t^2C$. Use Lemmas 2.10–2.12 to write $A^{(1)}$, B and C in terms of A . One finally gets the following equation which implies the result:

$$1 - A^{-1} = (t^4 + t^2)A + 2t^5A^2(1 - tA)^{-1} + t^6A^3(1 - tA)^{-1} + t^7A^4(1 - tA)^{-2}. \quad \square$$

Using MAPLE to solve the polynomial equation of Theorem 2.13, we obtain the following expression for $A = A(t)$:

$$\begin{aligned} & \{1 + 2t + (1 - 4t + 4t^2 - 4t^4)^{1/2} \\ & - \sqrt{2}(1 - 4t^2 - 2t^4 + (2t + 1)(1 - 4t + 4t^2 - 4t^4)^{1/2})^{1/2}\} / (4t^2), \end{aligned}$$

and a table of values (see Table 1).

Remark 2.14. There is an obvious bijection between A -paths of size n and words

Table 1

n	d_n	a_n	b_n	c_n	$a_n^{(1)}$
0	1	1	0	1	0
1	0	0	0	0	0
2	1	1	1	1	1
3	0	0	1	0	0
4	2	3	3	4	2
5	0	2	4	4	2
6	5	12	12	18	7
7	0	14	22	26	10
8	14	54	61	86	29
9	0	86	128	158	52
10	42	274	335	462	142
11	0	528	756	976	294
12	132	1515	1936	2665	772
13	0	3266	4580	6082	1732
14	429	8854	11652	16040	4451
15	0	20422	28402	38338	10482
16	1430	53786	72209	99536	26715
17	0	129368	179460	244880	64908
18	4862	336103	457274	631923	165194
19	0	830148	1151725	1583796	409720
20	16796	2145020	2945129	4081939	1044629

w in $xy^2, x^2y, x\bar{y}^2, x^2\bar{y}$ such that $\deg_x w = n, \deg_y w = \deg_{\bar{y}} w$ and $\deg_y v \geq \deg_{\bar{y}} v$ for $w = uv$; similarly for B -paths and C -paths.

3. Dyck paths of an (s, r) -knight

In Section 2, in order to find the generating series $A = A(t)$ for A -paths, we had to introduce B -paths (from $(0, 0)$ to $(n, 1)$) and C -paths (above $y = -1$). The generating series $A, B, C, A^{(1)}, B^{(0)}, C^{(0)}$ then satisfy the system of equations

- (1) $A = (1 - A^{(1)})^{-1}$,
- (2) $B = A \cdot B^{(0)}$,
- (3) $C = A + (B^{(0)})^2 \cdot A$,
- (4) $A^{(1)} = t^s A + 2t^s B + t^2 C$,
- (5) $B^{(0)} = t^2 A + tB$,
- (6) $C^{(0)} = A$.

For an (s, r) -knight, $s \leq r$ (which moves s squares in one direction and r squares perpendicularly) more new paths have to be introduced and a bigger system of equations has to be solved. For example, for a $(1, 3)$ -knight or a $(2, 3)$ -knight, we introduce for $i = 1$ and 2 , B_i -paths (from $(0, 0)$ to (n, i)) and C_i -paths (above $y = -i$). As before the integer k indicates the number of times the A -path or B_i -path (respectively the C_i -path) hits the line $y = 0$ (respectively $y = -i$).

Proposition 3.1. For a $(1,3)$ -knight the system of equations satisfied by the generating series $A, B_1, B_2, C_1, C_2, A^{(1)}, B_1^{(0)}, B_2^{(0)}, C_1^{(0)}$ and $C_2^{(0)}$ is

- (1) $A = (1 - A^{(1)})^{-1}$,
- (2) $B_1 = A \cdot B_1^{(0)}$,
- (3) $B_2 = A \cdot B_2^{(0)}$,
- (4) $C_1 = A + A \cdot (B_1^{(0)})^2$,
- (5) $C_2 = C_1 + A \cdot (B_2^{(0)})^2$,
- (6) $A^{(1)} = t^6 A + 2t^4 B_2 + t^2 C_2$,
- (7) $B_1^{(0)} = t^3 A + t B_2$,
- (8) $B_2^{(0)} = B_1^{(0)} \cdot B_1 + t B_1$,
- (9) $C_1^{(0)} = A$,
- (10) $C_2^{(0)} = C_1$.

Proof. Only (8) needs explanations: a $B_2^{(0)}$ -path either touches the line $y=1$ or it does not. \square

Theorem 3.2. The generating function for A -paths of a $(1,3)$ -knight satisfies the equation

$$t^8 A^8 - t^6 A^7 + (t^{10} - 2t^8 + t^6) A^6 + (2t^6 - t^4) A^5 + (2t^8 - 4t^6 + t^2) A^4 \\ + (2t^4 - t^2) A^3 + (t^6 - 2t^4 + t^2) A^2 - A + 1 = 0.$$

Proposition 3.3. For a $(2,3)$ -knight the system of equations is as in Proposition 3.1, except (6), (7) and (8) which are replaced by

- (6') $A^{(1)} = t^6 C_1 + 2t^5 (B_1 + B_1^{(0)} \cdot B_2) + t^4 C_2$,
- (7') $B_1^{(0)} = t^3 B_1 + t^2 B_2$,
- (8') $B_2^{(0)} = t^3 C_1 + t^2 (B_1 + B_2^{(0)} \cdot B_1)$.

Proof. For example, (6') is proved by erasing first and last steps in a $A^{(1)}$ -path. If first step is NNEEE and last step is SSSSE, then either it touches the line $y=1$ (giving the term $B_1^{(0)} \cdot B_2$) or it does not (giving the term B_1); the proof is similar for (8'). \square

Theorem 3.4. The generating function for A -paths of a $(2,3)$ -knight satisfies the equation

$$t^{16} A^8 - (2t^{15} + t^{12}) A^7 + (2t^{14} + t^{12} + 2t^{11}) A^6 - (2t^{13} + 2t^{11} + t^{10} + t^8) A^5 \\ + (t^{12} + 4t^{10} + t^4) A^4 - (2t^9 + 2t^7 + t^6 + t^4) A^3 \\ + (2t^6 + t^4 + 2t^3) A^2 - (2t^3 + 1) A + 1 = 0.$$

More generally for an (s,r) -knight, $s \leq r$, we have to introduce $r-1$ kinds of B -paths, that is B_i ($1 \leq i \leq r-1$) which goes from $(0,0)$ to (n,i) , and $r-1$ kinds of C -paths, that is C_i ($1 \leq i \leq r-1$) which stays above the line $y=-i$. The corresponding

$4r-2$ series $A, A^{(1)}, B_i, B_i^{(0)}, C_i, C_i^{(0)}$ ($1 \leq i \leq r-1$), where a $B_i^{(k)}$ -path (respectively $C_i^{(k)}$ -path) touches the line $y=0$ (respectively the line $y=i$) k times, satisfy the following theorem.

Theorem 3.5. *We have, using the natural notation $B_0=C_0=A$ and $B_{-1}=B_{-2}=\dots=0$,*

$$\begin{aligned}
 A &= (1 - A^{(1)})^{-1}, & A^{(1)} &= t^{2r}C_{s-1} + 2t^{r+s} \left(\sum_{1 \leq j \leq s} B_r^{(0)} \cdot B_{s-j} \right) + t^{2s}C_{r-1}, \\
 B_i &= A \cdot B_i^{(0)}, & B_i^{(0)} &= \sum_{1 \leq j \leq s} (t^r B_s^{(0)} \cdot B_{r-j} + t^s B_{r-j}^{(0)}) \cdot B_{i-j}, \\
 C_i &= C_{i-1} + A \cdot (B_i^{(0)})^2, & C_i^{(0)} &= C_{i-1} \quad \text{for } 1 \leq i \leq r-1.
 \end{aligned}$$

Proof. For example, we prove the fourth equation: given a $B_i^{(0)}$ -path, starting with the step $(0,0)$ to (r,s) , consider successively (for $j=1$ to s) if it touches the line $y=j$ or not. \dashv

Remark 3.6. Eliminating variables $C_i^{(0)}$ (trivially) and $B_i^{(0)}$ from these $4r-2$ equations leads to a system of $2r$ algebraic equations satisfied by the $2r$ generating series $A, A^{(1)}, B_i$ and C_i ($1 \leq i \leq r-1$). Further eliminations, as was done explicitly only in the cases of a $(1,1)$ -Dyck-knight (Theorem 1.9), a $(1,2)$ -chess-knight (Theorem 2.13), a $(1,3)$ -knight (Theorem 3.2) and a $(2,3)$ -knight (Theorem 3.4), give a polynomial equation (with polynomials in t as coefficients) of degrees 2, 4, 8 and 8 respectively satisfied by A .

Remark 3.7. Since the symmetry between s and r is never really used in Theorem 3.5, instead of (s,r) -, (r,s) -, $(s,-r)$ - and $(r,-s)$ -steps, one could as easily start with two (see [8] for the general case of m vectors) arbitrary vectors (s_1, r_1) and (s_2, r_2) in $\{1, 2, \dots\} \times \{0, 1, 2, \dots\}$, and consider Dyck paths made of (s_1, r_1) -, (s_2, r_2) -, $(s_1, -r_1)$ - and $(s_2, -r_2)$ -steps. A system of $4r-2$ equations, where $r = \max(r_1, r_2)$ similar to those above can be obtained.

Definition 3.8. Let $A_{s,r}(t)$ denote the generating series for A -paths of a (s,r) -knight.

Proposition 3.9. *We have $A_{s,r}(t^m) = A_{ms, mr}(t)$.*

Corollary 3.10. *The series $A_{m,m}(t)$ and $A_{m,2m}(t)$ have explicit closed forms.*

4. Conclusion

The case of several knights or several vectors can be studied similarly (see [8] for more general results using several variables series). One has to introduce more B -

paths and C -paths; the corresponding generating series satisfy an intricate system of algebraic equations. We may think of all these kinds of paths as words in the letters x, y, \bar{y} (see Remark 2.14); the corresponding languages are algebraic.

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